Strong-completeness and Faithfulness in Belief Networks

by

Chris Meek

January 1995

Report CMU-PHIL-62

Carnegie Mellon

Philosophy
Methodology
Logic

Pittsburgh, Pennsylvania 15213-3890

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Technical Report CMU-PHIL-62

January 27, 1995

[1]I would like to thank Clark Glymour, Thomas Richardson and Peter Spirtes for helpful conversations about the contents of this paper. Research for this paper was suported by the Office of Navel Research grant ONR #N00014-93-1-0568.

[2]Department of Philosophy, Carnegie Mellon University, Pittsburgh, PA 15213. E-mail address: cm1x@andrew.cmu.edu

1 Introduction

In a series of important papers, Geiger, Verma and Pearl (Geiger et al. 1990 and 1988 and Pearl 1988) outlined an axiomatic approach to characterizing and inferring independence relations in graphical statistical models. A class of graphical models of particular interest are the class of directed acyclic models called belief networks. Pearl introduced d-separation as a rule to infer the independence facts implied by a particular directed acyclic graph; an alternative equivalent rule has been proposed by Lauritzen et al. (1990). Geiger et al. (1990) have shown that d-separation is atomic-complete for independence statements for discrete Bayes networks; One is able to infer all of the atomic independence facts that can be legitimately inferred from the structure of the Bayes network. In this paper I show that d-separation has the property of strong-completeness for discrete Bayes networks; complete for arbitrary disjunctive and conjunctive combinations of independence statements. This result shows that d-separation as a rule of inference can not be improved upon.

The measure-theoretic approach to the proof of strong-completeness has important implications for one major approach to learning belief networks. Broadly speaking, there are two types of approaches to learning belief networks; the scoring approaches (Bayesian, Likelihood, MDL) and the independence approaches (PC, CI, SGS). The independence approaches have been shown to be asymptotically reliable assuming that the population distribution stands in a certain relationship to the structure to be learned. The distribution which stands in this relationship to the structure has been called by many names, e.g., faithful, stable; and the structure has been named a perfect map of such a distribution. I demonstrate that faithful multinomial distributions exist for every directed acyclic graph and every discrete state space. Furthurmore, in a specific measure-theoretic sense, there are many more faithful distribution than unfaithful distributions.

The new results in this paper are about discrete Bayes networks. Strong-completeness and the existence of faithful distributions has been shown previously for the Gaussian case. The discrete (multinomial) case is of special interest since many of the applications of machine learning and data modelling involve discrete data. I include the results for the Gaussian case and give a new and uniform proof of the results for both the Gaussian and multinomial cases.

2 Strong-completeness — the logic of belief networks

The basic goal of a logic is to derive statements entailed by the assumptions. In the case of the logic of belief networks we are interested in deriving independence statements from a directed acyclic graph $G = (V, E)$ which are true of any distribution V from a specific class of distributions over V for which G is an

I-map (see Pearl 1988)[1]. We use *V* to denote an arbitrary class of distributions, *Vjsf* to denote the class of multivariate normal distributions, and *Vv* for the class of multinomial distributions.

We let t range over independence statements, A J_L B|C for disjoint sets A, B and C and is read A is independent of B given C. We let I range over (i) independence statments and (ii) finite conjunctions and disjunctions of I statements. *G* entails / (written *G* f=p /) if and only if *I* is true in every distribution in *V* for which *G* is an I-map.

As with any logical calculi, there are rules of inference. The central rule of inference in this logic is that of d-separation. The first question one asks about a rule of inference is whether it is sound. The soundness of d-separation as a rule of inference has been demonstrated in Geiger et al. (1988). The next question one asks about a set of inference rule is whether the set of rules is complete; whether all of the true statements are derivable. A sentence / is derivable by a set of rules *V* from assumptions *G* (written *G \~v I)* if and only if there is a proof of/from *G* using the rules of inference *V*. Geiger et al. (1990) have shown that d-separation as a rule of inference is atomic-complete for the multinomial and multivariate normal class of distributions.

Theorem 4 [Geiger et al.] ***G \=Vv * */ an d on ty */ G ^v l-***

Theorem 3 [Geiger et al.] *G \=Vu* i *tf and only if G Vv* i-

Thus, for any given atomic independence fact i we can use d-separation to check if the independence statement i is entailed by the graphical structure. This does not allow us to check independence sentences which are disjunctive combination or sentences of disjunctive and conjuctive combinations of such statements. To derive disjunctive and conjunctive combinations of independence statements the set of inferential rules *V* must be expanded to a set of inferential rules *V** which includes A—introduction and V—introduction.

Theorem 11 [Strong-completeness] *G ^=V-D* I *if and only if G* hp* **I.**

Theorem 12 [Strong-completeness; Geiger et al., Spirtes et al.] *G \=Tjsr* I *if and only if G* hp* I.

The proofs of the strong completeness theorems are sketched in the appendix. As with any completeness proof, if a disjunctive independence sentence A _IL B|C V X JJ_ Y|Z is not true for a graph *G* we must show that there is a model — in our case a probability distribution — in which both A _IL B|C and X J_L Y|Z are not true. Let graph *G* be given and assume that it is not that case that *G \=-P-D* A _IL B|C. Geiger et al. (1990) gave a method for constructing a distribution *P* for which the given graph *G* is an I-map such A JJ_ B|C is false in P. I extend the result to show that there exists a distribution for arbitrary disjunctive combinations of non-entailed independence facts.

[1] Assuming that *P* has a density or probability mass function then *G* is an I-map for *P* if and only if *P* is Markov with respect to *G* (see Lauritzen et al. 1990 and Spirtes et al. 1993).

3 Assumptions for learning belief networks

There are several algorithms which use independence tests to learn belief networks from sample data including the PC, and SGS (Spirtes et al. 1993). Basically these algorithms perform a series of statistical tests of independence using the sample data and based upon the results of these tests the algorithms eliminate a set of possible models until the remaining set of models can not be distinguished by independence facts. The methods enumerated above differ in the series of independence test that are used; the selection and ordering of the tests can improve the practical reliability and computational tractability of these algorithms. Let S_G be any arbitrary boolean combination of independence statements about variables in graph G; we write S when the appropriate graph is clear from context. We interpret $\neg\ \mathbf{X} \perp\!\!\!\perp \mathbf{Y}|\mathbf{Z}$ to mean that $\mathbf{X}$, and $\mathbf{Y}$ are conditionally dependent on $\mathbf{Z}$. A distribution P is **faithful** to the graphical structure G if and only if exactly the independence facts true in P are entailed by the graphical structure G. We say that G **faithfully entails** S (written $G \models^{\mathcal{F}}_{\mathcal{P}} S$) if and only if S is true in every distribution in $\mathcal{P}$ which is faithful to G. It is easy to show that for all directed acyclic graphs G and for all S that $G \models^{\mathcal{F}}_{\mathcal{P}} S$ or $G \models^{\mathcal{F}}_{\mathcal{P}} \neg S$. Using this fact we can show the theoretical reliability of these algorithms assuming the correctness of the statistical tests and that the population distribution is faithful to the underlying graphical structure. Let $test_i$ be the result of the i^{th} test (e.g. $\mathbf{X} \perp\!\!\!\perp \mathbf{Y}|\mathbf{Z}$ or $\neg\mathbf{X} \perp\!\!\!\perp \mathbf{Y}|\mathbf{Z}$ for disjoint subsets $\mathbf{X}$,$\mathbf{Y}$,$\mathbf{Z}$ of vertices).[2] From the assumption, we can eliminate models in the following way. After performing t tests one can eliminate a model G if $\neg \bigwedge_{i=1}^{t} test_i$ is faithfully entailed by G. We can eliminate models until the remaining set of models are not destinguishable by conditional independence facts.

The assumption of faithfulness has been criticized by several researchers. The essence of the criticism is captured by the following question. How can one ever be confident that the population is faithful to the underlying structure?

This is a reasonable question but even stronger question seems warrented. Are there faithful distributions (in the class of distributions of interest) for any arbitrary directed acyclic graph? The theorems below demonstrate that the answer to this question is affirmative. The proof of existence for $\mathcal{P}_{\mathcal{N}}$ uses an alternative proof technique as compared to the proof given in Spirtes et al. (1993).

Theorem 9 [Existence] *For all directed acyclic graphs G there exists a $P \in \mathcal{P}_{\mathcal{D}}$ which is faithful to G.*

Theorem 10 [Existence—Geiger et al.] *For all directed acyclic graphs G*

[2] As above we can define a logical calculus for faithful derivability ($G \vdash_{\mathcal{F}} S$) using the rule of d-separation to derive both independence and dependence facts. By adding a complete set of propositional inference rules we can show that this logical calculus is strongly-complete for S sentences.

there exists a P £ Vjsf which is faithful to G.

But these theorems do not answer the criticism of the unreasonableness of the assumption of faithfulness. The next theorem shows that at least in a measure-theoretic sense the assumption of faithfulness is reasonable. The distributions in *Vv* and *Vjsf* are parametric distributions. Let TT^ be the set of linearly independent parameters needed to parameterize a discrete distribution for which graph *G* is an I-map and let *w^* be the set of linearly independent parameters needed to parameterize a multivariate normal distribution for which graph *G* is an I-map.

Theorem 7[Measure zero] *With respect to the Lebesgue measure over TT%, the set of distributions which are unfaithful to G is measure zero.*

Theorem 8 [Measure zero—Spirtes et al.] *With respect to the Lebesgue measure over 7r%, the set of distributions which are unfaithful to G is measure zero.*

4 Final remarks

I conjecture that the proof techniques presented in this paper can be extened to prove analogous measure zero, existence and strong-completeness results for the conditional Gaussian class of distributions (see Whittaker 1990).

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5 Appendix A — Proof sketches

The details given in this section are a bare-bones sketch of the proofs of the theorem in this paper. Detailed proofs can be found in the following section. Let G be some directed acyclic graph and π_G be the set of linearly independent parameters needed to encode any multinomial or multivariate normal distribution for which G is an I-map. As the context demands, we let π_G represent the parameters for either a multinomial or multivariate normal distribution.

Claim 1 *The independence facts not entailed by d-separation applied to directed acyclic graph G hold only for values of the parameters which statisfy non-trivial polynomial constraints.*

The proof of this claim is in two parts. First one can show, based upon the specific parameterization (multinomial or multivariate normal) that the constraints are polynomials in the parameters. Second we show that the constraints are non-trivial (not all value of the parameters satisfy the constraints). The proof of the non-triviality is similar to the main lemma used in the atomic-completeness result of Geiger et al..

Claim 2 *For independence statement* I *not entailed by G and for the Lebesgue measure over the set of parameters π_G the set of values where the independence fact* I *holds is Lebesgue measure zero.*

The proof of this claim follows from the fact that the solution set to non-trivial polynomial constraints has measure zero (See Spirtes 1994, unpublished).

Theorem 7 and Theorem 8 follow from Claim 2. With respect to a given graph G, only a finite number of independence facts are not faithfully entailed. Each of these independence facts hold only for a set of parameterizations of measure zero. The union of all of these finitely many sets of parameterizations is measurable and is of Lebesgue measure zero.

Theorem 9 and Theorem 10 follow from Theorem 7 and Theorem 8 by the following measure-theoretic argument. Given that the set of parameterizations in which the distribution is unfaithful are of measure zero and that there are sets of (permissable) parameterizations with positive measure then there are parameterizations which are faithful.

Finally, Theorem 11 and Theorem 12 and the strong-completeness for S sentences with respect to $\vdash_{\mathcal{F}}$ follow from the existence of faithful distributions for the two classes of distributions; all and only the idependence facts which follow from the rule of d-separation hold in the faithful distribution.

6 Appendix B — Proofs

A discrete Bayes network is a tuple $\langle G, P \rangle$ where P is a probability function over a finite set of variables $\mathbf{V}$ (each of which take on at least 2 values) and G is a graph over the same set of variables V such that there exists a factorization of *P{*V) such that

$$P(\mathbf{V}) = \text{II}_{A \in \mathbf{V}} P(A \backslash parents(A))$$

where *P(A\parents(A))* is a conditional probability distribution and *parents(A)* is the set of parents of vertex *A* in graph *G.* Linear Bayes networks are described in Spirtes et al. (1993).

6.1 Parameterizations of Bayes networks

Given that the joint distribution *P* factors according to the graph *G* into conditional probabilities we can parameterize the joint distribution by parameterizing each of the conditional distributions. For variable i E V w e define NV(A) as the number of possible values that *A* can take and C(^4) as the set of possible values of *A.* Let ˙sn*t(0) = 1 and *inst{{Al, ...,An})* = NV(J11) X ... X NV(^n). For each conditional probability *P(A\parents(A))* we can represent the conditional distribution with *nparam(A,parents(A))* linearly independent parameters where

$$nparam(A,parents(A)) = (\text{NV} - 1) \times inst(parents(A)).$$

The reason for the (NV(A) - 1) is that for any given instantiation of the parents the $NV(A)^{ih}$ probability is a linear combination of the other (NV(^4) - 1) parameters. We adopt the following convention for naming the parameters. *0D,d,(ai,...,aN)* is [tne] parameter for the conditional probability *P(D = d\Al =* a1, *...,AN = aN)* where *parents(D) = Al, ...,AN* and where the varibles (^41,..., *AN)* are ordered lexicographically. Let *0* be the set of all of the parameters for all of the variables in V. Each of the parameters *0D,d,(ai,...,aN)* satisfies the constraints that 0_< *0D,d,(ai,...,aN)*_< 1 and for all *al,...,aN* it is the case that *Y^d£C(D)^D,d,(ai,...,aN)* = I-[3] The parameterization of linear Bayes networks is discussed in Spirtes et al. (1993).

6.2 Faithfulness

Let *P* be a probability function (a density or mass function as required) over V and *G* be a graph over the vertices in V. *(G,P)* are said to satisfy the Markov condition if and only if for all *A* € V it is the case that *A* _IL *V\(parents(A)* U *descendants(A))* given the *parents(A).* If *P* is a discrete probability function then (G,P) is a discrete Bayes network if and only if *(G,P)* satisfies the Markov condition. If *P* is a multivariate normal density function then *(G,P)* is a Gaussian Bayes network if and only if *(G,P)* satisfies the Markov condition. A conditional independence relation is entailed for

[3] Note that if the distribution P is not positive then some of the parameters are not strictly necessary to parameterize the distribution.

graph G if and only if it is true in all distributions P such that $\langle G, P\rangle$ satisfies the Markov condition. Let $\langle G, P\rangle$ be a discrete (Gaussian) Bayes network. A discrete (Gaussian) Bayes network is faithful if and only if the conditional independence relations true in P are exactly those entailed by the factorization of P with respect to G.

6.3 Method for constructing constraints

In this section we give a method for calculating the polynomial constraint that must be satisfied for a violation of faithfulness to occur.

Let $\langle G, P\rangle$ satisfy the Markov condition and let $A \perp\!\!\!\perp B|\mathbf{C}$ be an independence fact true in P but such that $A \perp\!\!\!\perp B|\mathbf{C}$ is not implied by the Markov condition applied to G.

$$A \perp\!\!\!\perp B|\mathbf{C} \Leftrightarrow$$

$$\begin{gathered}\forall(a,b,\mathbf{c}) \text{ if } P(\mathbf{C}=\mathbf{c}) \neq 0 \text{ then} \\ P(A=a, B=b|\mathbf{C}=\mathbf{c}) = \\ P(A=a|\mathbf{C}=\mathbf{c})P(B=b|\mathbf{C}=\mathbf{c}) \Leftrightarrow\end{gathered} \tag{$*$}$$

$$\begin{gathered}\forall(a,b,\mathbf{c}) \text{ if } P(\mathbf{C}=\mathbf{c}) \neq 0 \text{ then} \\ P(A=a, B=b, \mathbf{C}=\mathbf{c})P(\mathbf{C}=\mathbf{c}) = \\ P(A=a, \mathbf{C}=\mathbf{c})P(B=b, \mathbf{C}=\mathbf{c})\end{gathered}$$

Thus, for a violation of faithfulness to occur a set of $inst(\{A,B\} \cup \mathbf{C})$ many equations must be satisfied. Now I will show that each of these equations is a polynomial in the parameters of the model. Let $\Omega = Ancestor(A, B, \mathbf{C})$.

$$P(\Omega) = \textstyle\sum_{\mathbf{V}\setminus\Omega} P(\mathbf{V}) = \prod_{X\in\Omega} P(X|parents(X)) \tag{1}$$

Notice that all of the probabilities which occur in the last statement in (*) are of the form $P(\mathbf{F}{=}\mathbf{f})$ for some set of vertices $\mathbf{F}$ and some instantiations of those vertices $\mathbf{f}$. For instance if $\mathbf{F} = \{A\} \cup \{B\} \cup \mathbf{C}$ then $P(\text{A=a,B=b, }\mathbf{C}{=}\mathbf{c})$ is of the form $P(\mathbf{F}{=}\mathbf{f})$.

$$\begin{aligned} P(\mathbf{F}=\mathbf{f}) &= \textstyle\sum_{\Omega\setminus\mathbf{F}} P(\mathbf{F}=\mathbf{f}, \Omega\setminus\mathbf{F}) && (2)\\ &= \textstyle\sum S(i) && (3)\end{aligned}$$

where each summand $S(i)$ is a product of $|\Omega|$ many parameters (see equation 1). There are $I = inst(\Omega\setminus\mathbf{F})$ many summands. Let $O1, \ldots, OI$ be the I instantiations of the variables in $\Omega\setminus\mathbf{F}$. We define $M(A,i)$ as follows

$$\begin{aligned} M(A,i) &= \text{the value of } A \text{ in the } i^{th} \text{ instantiation of } \Omega\setminus\mathbf{F} \text{ if } A \in \Omega\setminus\mathbf{F}\\ &= \text{the value of } A \text{ in the instantiation } \mathbf{f} \text{ if } A \in \mathbf{F}.\end{aligned}$$

The i^{th} summand $S(i)$ for (2) is formed in the following fashion. Let $p(A,j)$ be the j^{th} parent of A with respect to the lexicographic ordering over the set of parents of A.

$S(i) = \prod_{A \in \Omega} \theta_{A,M(A,i),(M(p(A,1),i),\ldots,M(p(A,n),i)}$

where A has n parents

Since the constraints are products of terms of the form $P(\mathbf{F}=\mathbf{f})$ given in (2) the constraints are polynomials in the parameters of the discrete Bayesian network.

Violations of faithfulness occur in the linear case for distributions only if polynomial constraints in the parameters of the model hold. This is shown in Spirtes et al. (1993).

6.4 The polynomial constraints are non-trivial

A polynomial in n variables is said to be non-trivial (not an identity) if not all instantiations of the n variables are solutions of the polynomial. Now we show that all of the polynomials for non-entailed independence constraints are non-trivial.

We do this by using the property weak transitivity which is guaranteed to hold in Gaussian and discrete distributions where all of the variables are binary (see Pearl 1988). Weak transitivity allows us to give an alternative proof of the completeness of d-separation and a measure theoretic result about faithfulness for the Gaussian case as well as the discrete case.

Some inference rules about independence and dependence for probability theory (see Dawid 1979 and Pearl 1988).

$\neg X \perp\!\!\!\perp Y|Z \Rightarrow \neg Y \perp\!\!\!\perp X|Z$ (Symmetry)

$\neg X \perp\!\!\!\perp Y|Z \Rightarrow \neg X \perp\!\!\!\perp WY|Z$ (Decomposition)

For positive distributions

$\neg X \perp\!\!\!\perp WY|Z \wedge X \perp\!\!\!\perp W|ZY \Rightarrow \neg X \perp\!\!\!\perp Y|ZW$ (Intersection)

The following rule also holds for Gaussian and Boolean systems

$\neg X \perp\!\!\!\perp \gamma|Z \wedge \neg\gamma \perp\!\!\!\perp Y|Z \Rightarrow \neg X \perp\!\!\!\perp Y|Z \vee \neg X \perp\!\!\!\perp Y|Z\gamma$ (Weak Transitivity)

where γ is a singleton set.

Lemma 1 — *If in directed acyclic graph G there exists a d-connecting path between A and B given $\mathbf{C}$ then there exists a singly-connected subgraph G' of G such that A and B given $\mathbf{C}$ are d-connected by a path p in G' and such that the only edges in G' are edges on the d-connecting path and a set of edges which form exactly one directed path from each collider on the path p to a member of $\mathbf{C}$.*

Proof — Let p be a d-connecting path between A and B given $\mathbf{C}$ in G. Let $G1$ be the subgraph of G such that all of the edges on p are in G_1 and for each collider D on p not in $\mathbf{C}$ we include the edges that are on one path from D to a member of $\mathbf{C}$ not through another member of $\mathbf{C}$. Arbitrarily choose one path if there are more than one from D to members of $\mathbf{C}$.

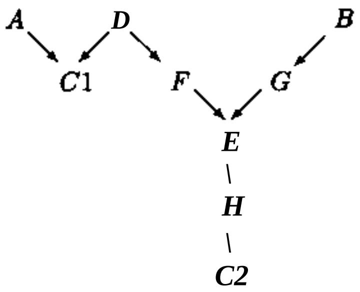

Figure 1: Schematic of singly connected graph between *A* and *B* given C = {C1,C2}.

Let r(Gi) be the number of multiple pathways that exist in Gl. It is clearly finite. If r(Gi) > 0 then there exists two distinct colliders onp, the d-connecting path from *A* to *B* given C in Gi, which have the same member of C as a descendent. Let *D* and *D2* be such colliders. Let *p[Dl,D2]* be the set of edges between *D* and *D2* on the path *p.* Remove *p[Dl,D2]* from the graph *Gl.* Clearly we reduce r(C?i) by removing *p[Dl, D2],* Continue the process until r(Gi) = 0. Clearly a d-connecting path between *A* and *B* given C remains at each stage.

The graph *Gi* is the desired graph (?'.•

Schematically, the claim amounts to the claim that there is a subgraph which looks like the graph in Figure 1 where C = {Cl, C2}.

Let *P* be a probability distribution and *G* be a directed acyclic graph. (G, *P)* satisfies the *local dependence* condition if and only if (G, *P)* is a Bayes network and if *A* —+ *B* is in *G* then -A* _IL *B* is true in *P*.**

Claim 3 *For a singly-connected graph G there exists a positive binary probability distribution P distribution such that (G,P) satisfy the local dependence condition.*

Claim 4 *For a singly-connected graph G there exists a positive Gaussian probability distribution P distribution such that* (G,P) *satisfy the local dependence condition.*

Both Claim 3 and Claim 4 are easy to show.

Lemma 2 *If (G,P) satisfies the local dependence condition, P is weakly transitive, G is singly-connected, and there exists a directed path from Ai to A_n in G then ^A\ JJ_ A_n is true in P.*

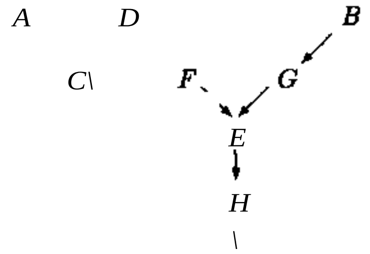

Figure 2: Graph *G'*.

Proof— by induction on length of path using weak transitivity.

base case — consider the trivial case of the null path from *A* to *A*.

induction step — Assume that there is a path from *A* to A_n in *G* and *->A* JJL *An-*. We get *->An-i* JJL *An* from local dependence and *Ai* i_L A_n \\$A_n$*-i* by the Markov condition and the single-connectedness of *G*. Then, by weak transitivity, $\neg A_1 \perp\!\!\!\perp A_n$.□

Theorem 3 *For any directed acyclic graph G which does not entail A* _ÍL *B\C there exists a discrete binary (Gaussian) distribution P such that ->A* _ÍL *B\C is true in P and* (G,P) *is a discrete (Gaussian) Bayes network.*

Proof— assume that *A* U_ *B\C* is not entailed by Markov condition applied to *G*. We construct a binary valued (Gaussian) distribution Pi over the variables in *G* such that *(G,P)* is a discrete Bayes network and such that *-iA* JJ_ *B\C* is true in Pi. Since *A* JJL *B\C* is not entailed by Markov condition applied to *G* there must exist a path which d-connects *A* and *B* given C. Let *G'* be the subgraph of *G* described in Lemma 1. To simplify the proof I will give an informal argument which can readily be turned into a rigorous inductive argument. Let *G'* be described by the graph in Figure 2 where C = {Cl, *C2}*.

Let P be a positive binary (Gaussian) probability distribution such that *{G'jP)* satisfies the local dependence condition; one exists by Claim 3 (Claim4). Note that the positivity of P allows us to use Intersection as a rule of inference.

The goal is to show that -^*A 1± B\C1, C2*

(1) -.A JJL *D\C1*

proof —

(i) *A 1± D*	Markov condition (applied to G^1)
(ii) -.£> JJ_ C1|0	local dependence
(iii) -<C1 JJ_J4|0	local dependence

(iv) -.A 11 *D\C* weak transitivity (WT) and (i), (ii) and (iii)

(2) $\neg A \perp\!\!\!\perp D|C1, C2$

proof -

(i) -u4 JJL *D, C2\C1* from Decomposition and (1)

(ii) *A* 11 *C2\D,* Cl from Markov condition

(iii) -**A* 11 *D\C1, C2* from (i), (ii) and Intersection

(3) -•£> 1 1 *B\C2*

proof-

(i) *D* 11 £|0 from Markov condition

(ii) -*D* 11 C2|0 local dependence and Lemma 2

(hi) -iC2 11 *B\$* local dependence and Lemma 2

(iv) -.£> 11 *B\9* or -i£> 11 £|C2 from WT, (ii) and (iii)

(v) -.£> 11 *B\C2* from (i) and (iv)

(4)-.£>JLL£|Cl,C2

proof- as in proof of (2).

(5)-«Al±B\Cl,C2

proof-

(i) *A* 11 Z?|C1, C2, Z) from Markov condition

(ii) -.i4 11 £|C1, C2 or -^4 11 £|C1, C2, Z? from WT, (4) and (2)

(iii) -A 11 B|C1, C2 from (i) and (ii)

Thus we have established that -**A* 11 *B\C\,C2* and it is clear that we can extend *P* to a distribution *P* over G. Let V^7 be the set of vertices in *G'* and V be the set of vertices in *G* and let *{Zi, Z2, >>*Z_n*}* be an enumeration of the vertices in V\V. In the discrete case let *P(V)* = Pi(V')P(^i)... *P(*Z_n*)* where *P(Z{)* is any arbitrary binary distribution over the variable *Zi.* In the Gaussian case let cov(Z_t-, *X)* = 0 for all $0 \leq i \leq n$ and *X* E V and set the variances of *Zi* arbitrarily but not to zero. It should be clear that the proof above can be turned into an induction over the number of directed paths (or more exactly semi-treks) in the d-connecting path in G'.D

Theorem 4 (Geiger et al.; Atomic completeness) *For any directed acyclic graph G over variables* V *which does not entail* *A* 11 *B\C there exists a discrete (not necessarily binary) distribution P such that -<A* 11 *B\C is true in P and (GjP) is a discrete Bayes network.*

Proof — Begin by constructing the discrete binary distribution *P* from Theorem 3. We must simply expand the distribution based on binary valued probabilities to one based upon the number of categories required for each of the variables in V; the resulting distribution is essentially a binary distribution extended to an arbitrary discrete probability space by using zero probabilities. We assume that the values of the binary variables in *P* are either zero or one (0 or 1). The easiest way to extend the distribution is to force the probability of $\mathbf{V}=\mathbf{v}$ to be zero if for some $A \in \mathbf{B}$ the value of A in $\mathbf{v}$ is not either 0 or 1. The dependence follows since all of the polynomials described in equation (*) must hold for the independence to hold and by Theorem 3 and this is not the case.□

6.5 Polynomial constraints are Lebesgue measure zero and the completeness of d-separation.

Theorem 5 (Spirtes unpublished) - *The solutions to a (non-trivial) polynomial are Lebesgue measure zero over the space of the parameters of the polynomial.*

For a fixed statespace (i.e. the number of categories for each variable) let $\pi_G^{\mathcal{D}}$ be the set of linearly independent parameters needed to parameterize an arbitrary discrete distribution for which graph G is an I-map and let $\pi_G^{\mathcal{N}}$ be the set of linearly independent parameters needed to parameterize an arbitrary multivariate normal distribution for which graph G is an I-map. For the discrete case, the set of legal parameterizations $\mathrm{E} \subseteq [0,1]^n$ where n is the number of linearly independent parameters. For the Gaussian case, the set of legal parameterizations is the space $\mathcal{R}^n$.

Theorem 6 *For a fixed statespace with n linearly independent parameters, the set of parameterizations ω over a graph G in which independence fact $A \perp\!\!\!\perp B|\mathbf{C}$ is true but such that $A \perp\!\!\!\perp B|\mathbf{C}$ has measure zero with respect the Lebesgue measure over $\mathcal{R}^n$.*

Proof — Let $n = \mathit{inst}(\{A, B\} \cup \mathbf{C})$. There are n polynomials which must hold for this violation to occur. The polynomials are non-trivial by Theorem 3. Let ω_i be the set of solutions to the i^{th} polynomial.

$$\omega = \bigcap_{i=1}^{n} \omega_i$$

ω is measurable since finite intersections of measurable sets are measurable. Let $\omega' = \cup_{i=1}^{n} \omega_i$. Since ω' is the finite union of measurable sets it is measurable. $\mu(\omega) \leq \mu(\omega') \leq \sum_{i=1}^{n} \mu(\omega_i) = 0$ and given the non-negativity of the measure we have $\mu(\omega) = 0$; Ω is Lebesgue measure zero.□

Theorem 7 *Violations of faithfulness in a discrete bayes network $\langle G, P\rangle$ are measure zero with respect to the Lebesgue measure over $\mathcal{R}^n$ where $|\pi_G^{\mathcal{D}}| = n$.*

Proof — There are a finite number of sets of polynomials which must be satisfied to violate faithfulness. Let n be the number of such polynomials and ω_i be the set of solutions to the i^{th} set of polynomials, each of these sets is of measure zero by Theorem 6.

$$\omega = \bigcup_{i=1}^{n} \omega_i$$

ω is measurable since finite unions of measurable sets are measurable and the set ω is also Lebesgue measure zero. Finally restrict the solution set ω to the interval [0,1]. Let $\mathrm{E} \subseteq [0,1]^n$ be the subset of legal parameterizations of a distribution where n is the dimensionality of the space for the Lebesgue measure, μ. E is a closed set and thus measureable. As E is a measurable set and the finite intersection of measurable sets is again measurable, $\omega \cap \mathrm{E}$ is measurable. Since $\omega \cap \mathrm{E} \subset \omega$ we know that $\mu(\omega \cap \mathrm{E}) \leq \mu(\omega) = 0$ and by the non-negativity of the measure μ we have that $\mu(\omega \cap \mathrm{E}) = 0$.□

Theorem 8 *Violations of faithfulness in linear probability distributions are Lebesgue measure zero.*

Proof — Similar to proof of Theorem 7.□

Theorem 9 *For all directed acyclic graphs G there exists a $P \in \mathcal{P}_{\mathcal{D}}$ which is faithful to G.*

Proof — Follows from Theorem 7 by the following measure-theoretic argument. Given that the set of parameterizations in which the distribution is unfaithful are of measure zero and that there are sets of (permissable) parameterizations with positive measure then there are parameterizations which are faithful.□

Theorem 10 (Existence—Spirtes et al.) *For all directed acyclic graphs G there exists a $P \in \mathcal{P}_{\mathcal{N}}$ which is faithful to G.*

Proof — Follows from Theorem 8 as in Theorem 9.□

Theorem 11 *d-separation is complete for the class of multinomial distributions over arbitrary directed acyclic graphs.*

Proof — This follows from the existence of a faithful distribution.□

Theorem 12 *d-separation is complete for the class of Gaussian probability distributions over arbitrary directed acyclic graphs.*

Proof — This follows from the existence of a faithful distribution.□